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#### Armored

#### CPU-bound Encryption for Android-driven ARM Devices

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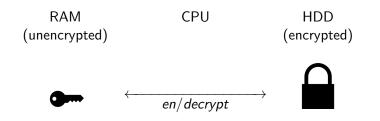
# Full Disk Encryption

- Full disk encryption (FDE) protects data against *physical lost* and theft of the hard drive
- It does not protect against remote attacks



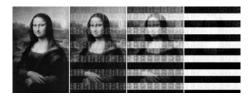
#### Software Disk Encryption

Current (software-based) FDE solutions do *not* protect data effectively if an adversary gains *physical access*!



### Coldboot Attack

# Disk Encryption Key in RAM $\rightarrow$ Exploit remanence effect of RAM





# Encryption on Android

#### Encryption on Android

- Since Android 4.0 aka Ice Cream Sandwich (ICS)
- Based on dm-crypt (device-mapper and Linux Crypto API)
- Only the user partition /data is encrypted
- Mode aes-cbc-essiv:sha256 is enforced with 128-bit keys

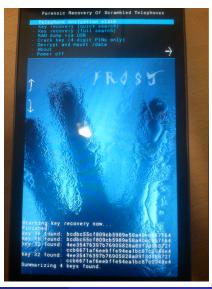
Of course encryption is possible with all common Linux distributions that run on ARM, too!

Encryption on Android

#### Coldboot Attack with Smartphones FROST: Forensic Recovery Of Scrambled Telephones

And it works with smartphones too! In this example: Galaxy Nexus





Johannes Götzfried, Tilo Müller (FAU)

Memory attacks require target systems to be *running* or *suspended*:

- Lost and theft of suspended laptops
- Confiscation of running servers
- But smartphones are always on

Basically *all* memory contents can be read out  $\rightarrow$  We focus on the security of **disk encryption keys**!

# ARMORED Security Policy

On  $\boldsymbol{\mathsf{ARM}}$  we  $\boldsymbol{\mathsf{O}}\mathsf{bstruct}$  the Recovery of  $\boldsymbol{\mathsf{E}}\mathsf{ncryption}$  Keys from  $\boldsymbol{\mathsf{D}}\mathsf{RAM}$ :

- AES implementation solely on the ARM microprocessor
- No sensitive information enters RAM
  - secret keys
  - key schedule
  - all intermediate states
- Only processor registers are used as storage



#### It has already been done ...

... on x86:

- FrozenCache
- LoopAmnesia
- TRESOR
- TreVisor

#### TRESOR

- uses the x86 debug registers *dr0* to *dr3* as key storage
- utilizes SSE registers to execute the AES algorithm
- implements AES using AES-NI

But ARMORED is the first CPU-bound encryption for ARM devices!

# Challenge

#### Security Policy

No valuable information about the AES key or state should be visible in RAM at any time

 $\rightarrow$  Implement AES without using RAM at all

- no runtime variables
- no stack
- no heap

 $\rightarrow$   $\mathrm{Armored}$  core is written in pure ARM assembler

 $\rightarrow$  Misuse registers as key storage

# Key Storage

Mix of breakpoint and watchpoint registers:

- Only accessible from kernel space
- seldom used by end-users

#### Memory alignment

- instructions are 4 bytes long and 4 bytes aligned
- two least significant bits of break- and watchpoint registers are zero

ightarrow divide key-sequence into 16 bit chunks (for simplicity)

- store parts to the 16 most significant bits of the registers
- 4 breakpoint and 4 watchpoint registers: 128 bit
- PandaBoard: even 6 breakpoint and 4 watchpoint registers

 $\rightarrow$  AES-128 is possible, enough for Android's disk encryption

### Working Register Set

NEON register set as temporary working storage:

- SIMD instruction set
- supported by many chips, e.g. Cortex-A9
- sixteen 128-bit registers, i.e. 256 bytes
- 64-bit and 128-bit SIMD operations
- access on byte granularity

#### Example

/* regi	ister	defs	*/
rstate	.qn	q0	
rhelp	.qn	q1	
rk1	.qn	q2	
rk1d0	.dn	d4	
rk1d1	.dn	d5	

```
/* xor sbox(key[index]) onto r2 */
.macro ks_box index base rk
  vmov.u8 r3, \rk\()d0[\index]
  ldr r3, [\base, r3, lsl #2]
  eor r2, r2, r3
.endm
```

### Gladman's AES Method

TRESOR implementation relies heavily on AES-NI

• AES-128 consists of basically 10 times aesenc

ARM has no AES-NI instruction set

- $\rightarrow$  use Gladman's AES Method
  - based on table lookups
  - efficient without special hardware

#### Specialities with ARM assembler

- RISC: all instructions are 4 bytes
- 4-byte base address of table cannot be loaded as immediate value
- manually generate constant pool and store pool address to register
- get base address register indirect

# Key Schedule

Conventional AES:

• round keys are calculated *once* and then stored in RAM (for performance reasons)

Armored:

 on-the-fly round key generation (entire key schedule is too big to be stored inside the CPU)

.macro key_schedule       generate_rk       rk2, rk3         eor       r1, r1, r1       generate_rk       rk3, rk4         ldr       r7, [r12]       generate_rk       rk4, rk5         add       r8, r7, #1024       generate_rk       rk5, rk6         add       r9, r8, #1024       generate_rk       rk6, rk7         add       r10, r9, #1024       generate_rk       rk7, rk8         ldr       r11, [r12, #4]       generate_rk       rk8, rk9         generate_rk       rk1, rk1       generate_rk       rk9, rk10	Example			
generate_ik iki, ikz .enum	eor	r1, r1, r1	generate_rk	rk3, rk4
	ldr	r7, [r12]	generate_rk	rk4, rk5
	add	r8, r7, #1024	generate_rk	rk5, rk6
	add	r9, r8, #1024	generate_rk	rk6, rk7
	add	r10, r9, #1024	generate_rk	rk7, rk8
	ldr	r11, [r12, #4]	generate_rk	rk8, rk9

### Kernel Patch

 $\ensuremath{\operatorname{ARMORED}}$  is designed as a Linux kernel patch for three reasons:

- Interpret and Android FDE uses the Linux Crypto API
- Problem: swapping of registers due to context switches
   → Solution: introduce *atomicity*

ARMORED is implemented in kernel space (currently Linux 3.2)

pprox 1700 LOC pprox 500 lines assembly code



#### Atomic Sections

- OS regularly performs context switches
- if ARMORED is active this context comprises sensitive data  $\rightarrow$  run ARMORED atomically (per 128-bit input block)

#### Example

```
void encrypt(struct crypto_tfm *tfm,
             u8 *dst, const u8 *src)
ſ
  unsigned long irq_flags;
  preempt_disable();
  local_irq_save(irq_flags);
  encblk_128(dst, src);
  local_irq_restore(irq_flags);
  preempt_enable();
}
```

#### **Development Platform**

Main development and testing was done on a PandaBoard running with Ubuntu 12.04 LTS (Precise Pangolin)





A Galaxy Nexus running Android 4.0 (Ice Cream Sandwich) has been tested as well  $\operatorname{ArmORED}$ : nothing but the output block is written actively to RAM

But: sensitive data may be copied into RAM *passively* by OS side effects (e.g. interrupt handling, scheduling, swapping, etc.)  $\rightarrow$  observe RAM of a ARMORED system at runtime

#### Tests

- Use FROST to actually perform a coldboot attack
- Look for keyschedule in RAM using AESKeyFind
- Look for the key in RAM (search for longest match)

Physical RAM was dumped using LiME – the Linux Memory Extractor  $\rightarrow$  We did not find the key (longest match was 4 bytes)

How to ensure that our implementation is correct?

- Linux kernel provides a test manager
  - check with official AES test vectors
- Encrypt random data with ARMORED and decrypt with generic AES
- Encrypt random data with generic AES and decrypt with ARMORED

 $\rightarrow$  We have good evidence that our implementation is correct

At first  $\operatorname{Armored}$  was 4.5 times slower than generic AES

#### Improvement: Larger atomic sections

- Process more input blocks per atomic section
  - reduce number of necessary key schedules
- How many blocks per section?
  - interactivity is no problem (1-2 microseconds vs. milliseconds)
  - could make sections large (up to 1024 blocks)
  - but: only 512 bytes per sector, i.e. maximal 32 blocks
- Necessary to change modes of operation: ECB, CBC, CTR

 $\rightarrow$  two  $\operatorname{ARMORED}$  variants: 16 blocks or only 1 block per section

### Performance Results

Reading 400 MB random data from encrypted RAM disk:

- Generic AES: 15.55 MB/s
- Armored 1x: 3.57 MB/s
- Armored 16x: 6.76 MB/s

Comparison of coldboot resistant implementations:

	slowdown
TRESOR	1.5
TreVisor	1.5
LoopAmnesia	2.0
Armored 1x	4.5
Armored 16x	2.3

Installation of ARMORED on smartphones is not very easy

- A kernel patch is not a user friendly application
- You might even do not have the code or parts of it

Bootstrapping problems

- How to get the key into the debug registers?
- Currently via adb and a sysfs interface

Integration into the android boot prompt

- Would be easily possible
- Just change hardcoded cipher and use sysfs interface

#### Confidentiality

Almost impossible to ensure that no password or key fragments remain within RAM

# Conclusion

 $\ensuremath{\operatorname{ARMORED}}$  withstands coldboot attacks and protects your DEK lt does not prevent:

- Local privilege escalation
- JTAG attacks
- Loss of other sensitive data in RAM

 $\operatorname{ARMORED}$  can be used

- practical on ARM based laptops
- on smartphones only as proof of concept



 $\operatorname{Armored}$  is the first CPU-bound encryption for ARM devices

Thank you for your attention!

Further Information:

http://www1.cs.fau.de/armored

